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# Implementing Object-Oriented Languages - Part 2

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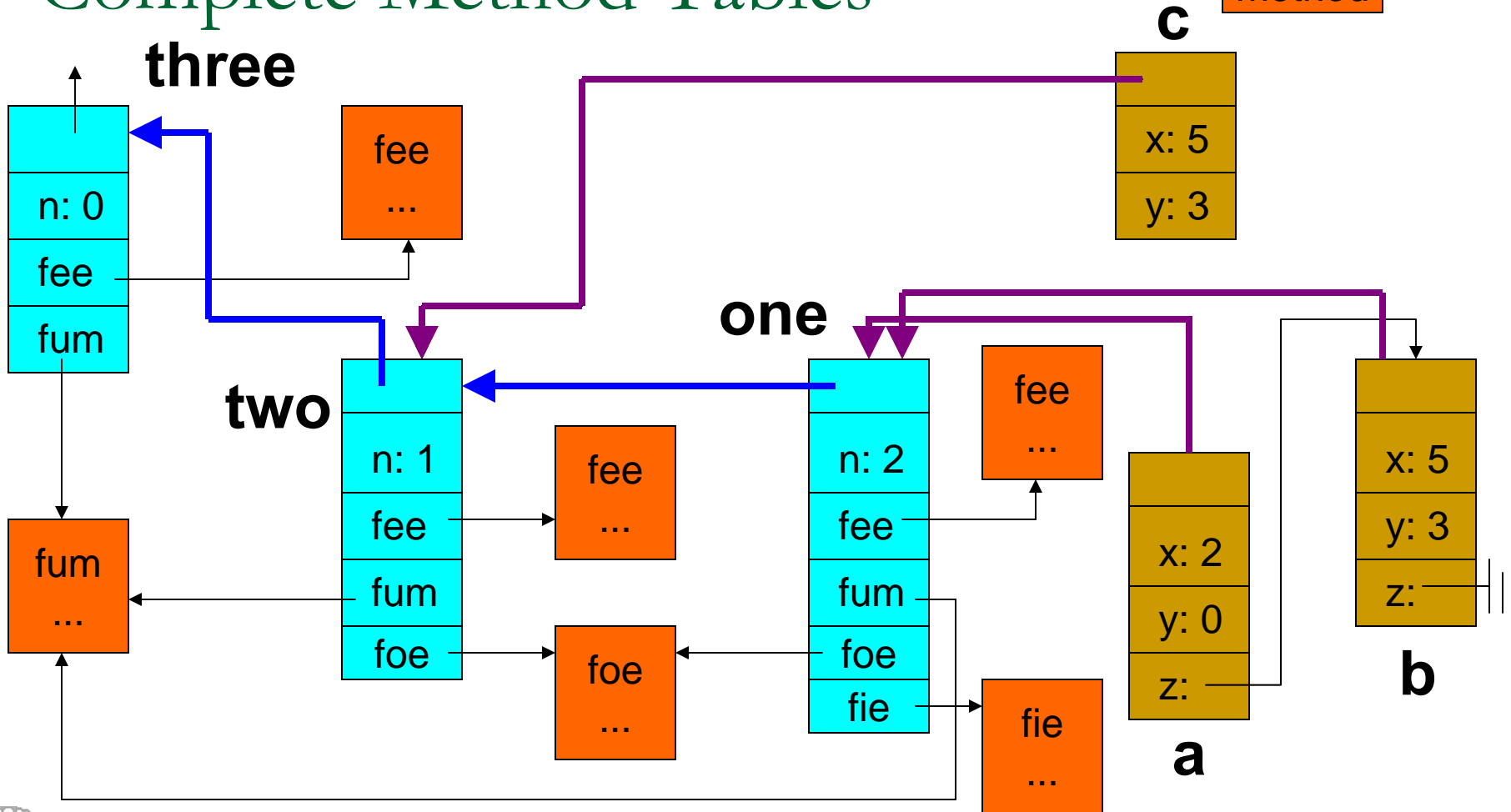
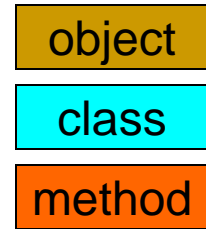
# Outline of the Lecture

- Language requirements
- Mapping names to methods
- Variable name visibility
- Code generation for methods
- Simple optimizations
- **Parts of this lecture are based on the book, “Engineering a Compiler”, by Keith Cooper and Linda Torczon, Morgan Kaufmann, 2004, sections 6.3.3 and 7.10.**

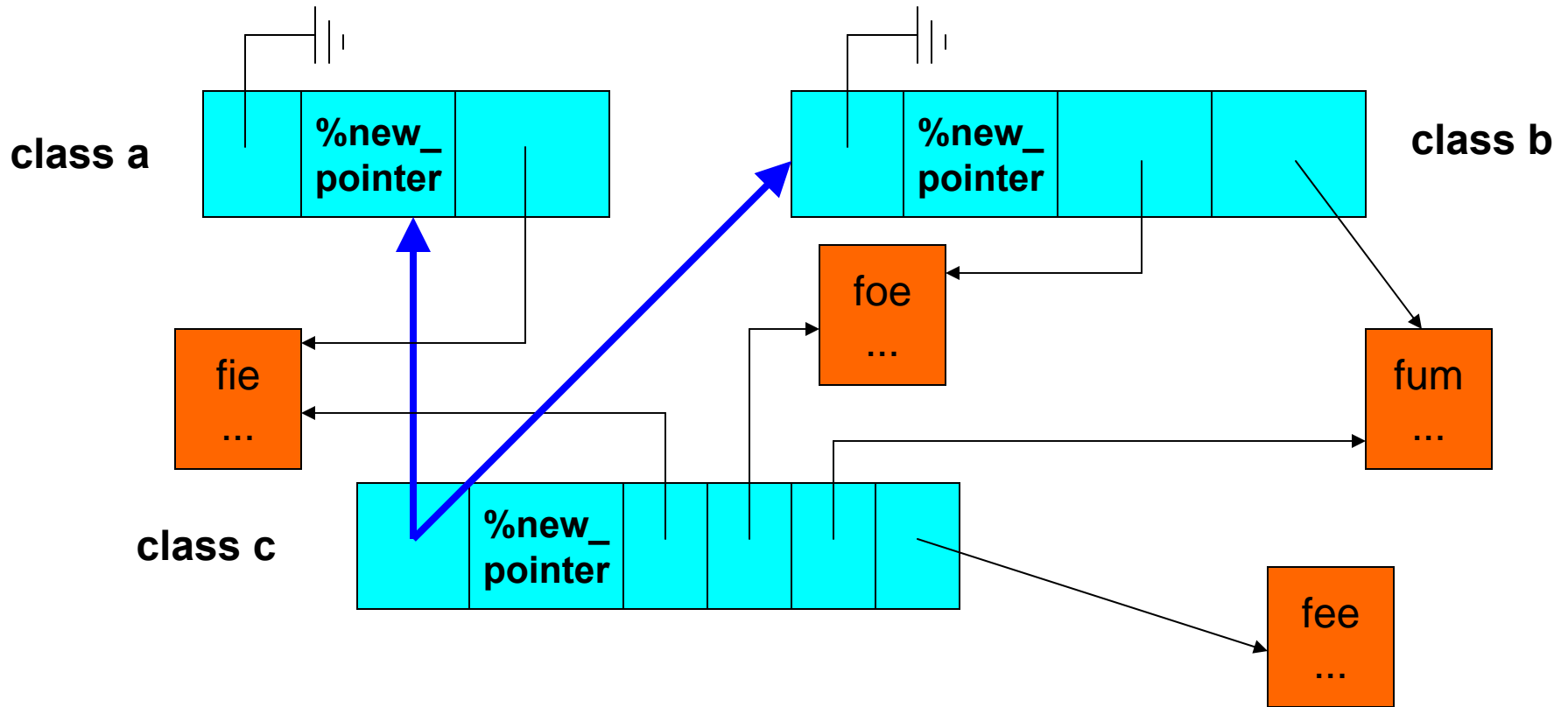
Topics 1,2,3, and, 4 were covered in Part 1 of the lecture.



# Example of Class Hierarchy with Complete Method Tables



# Implementing Multiple Inheritance



# Implementing Multiple Inheritance

object layout  
for objects  
of class a

<b>class pointer</b>	<b><i>a</i> data members</b>
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object layout  
for objects  
of class b

<b>class pointer</b>	<b><i>b</i> data members</b>
--------------------------	----------------------------------

object layout  
for objects  
of class c

<b>class pointer</b>	<b><i>a</i> data members</b>	<b><i>b</i> data members</b>	<b><i>c</i> data members</b>
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# Implementing Multiple Inheritance

## - Fixed Offset Method

object layout  
for objects  
of class c

<b>class pointer</b>	<b>a data members</b>	<b>b data members</b>	<b>c data members</b>
--------------------------	---------------------------	---------------------------	---------------------------

- Record the constant offset in the method table along with the methods
  - Offsets for this example are as follows:
    - (c) fee : 0, (a) fie: 0, (b) foe : 8, (b) fum : 8, assuming that instance variables of class a take 8 bytes
  - Generated code adds this offset to the receiver's pointer address before invoking the method



# Implementing Multiple Inheritance

## - Trampoline Functions

- Create **trampoline** functions for each method of class **b**
  - A function that increments **this** (pointer to receiver) by the required offset and then invokes the actual method from **b**.
  - On return, it decrements the receiver pointer, if it was passed by reference



# Implementing Multiple Inheritance

- Trampolines appear to be more expensive than the fixed offset method, but not really so
  - They are used only for calls to methods inherited from **b**
    - In the other method, offset (possibly 0) was added for all calls
  - Method inlining will make it better than option 1, since the offset is a constant
- Finally, a duplicate class pointer (pointing to class **c**) may need to be inserted just before instance variables of **b** (for convenience)



# Fast Type Inclusion Tests – The need

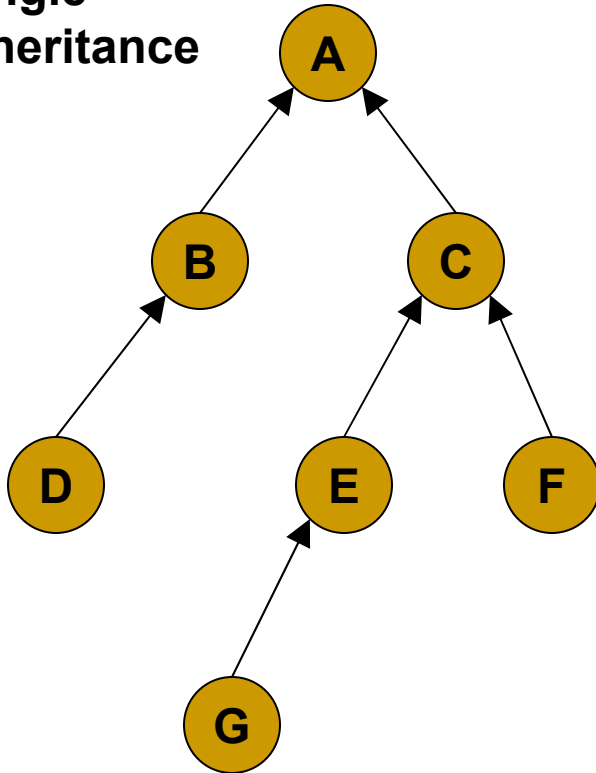
- If class **Y** is a subclass of class **X**
  - `X a = new Y();` //a is of type base class of Y, okay  
// other code omitted  
`Y b = a;` // a holds a value of type Y
  - The above assignment is valid, but the following is not
  - `X a = new X();`  
// other code omitted  
`Y b = a;` // a holds a value of type X
- Runtime type checking to verify the above is needed
- Java has an explicit **instanceof** test that requires a runtime type checking

# Fast Type Inclusion Tests – Searching the Class Hierarchy Graph

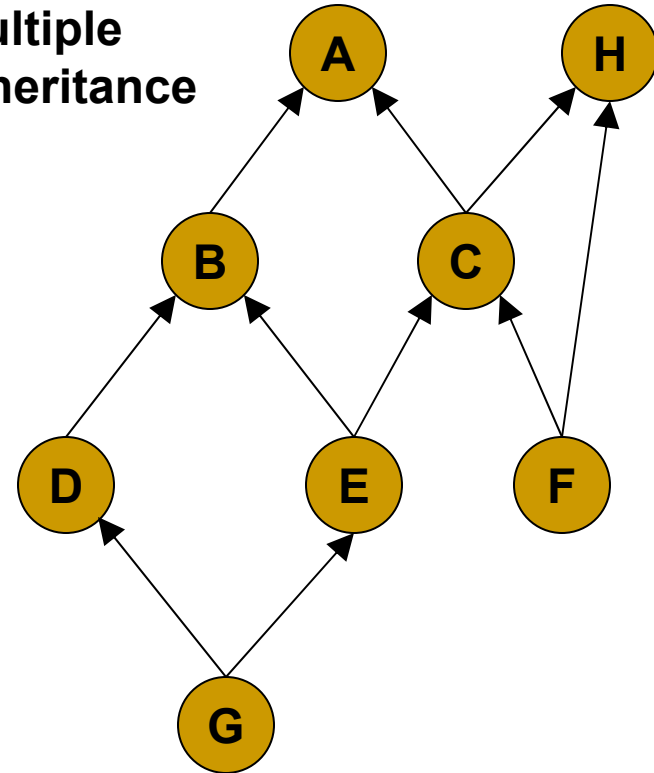
- Store the class hierarchy graph in memory
- Search and check if one node is an ancestor of another
- Traversal is straight forward to implement only for single inheritance
- Cumbersome and slow for multiple inheritance
- Execution time increases with depth of class hierarchy

# Class Hierarchy Graph - Example

Single inheritance



Multiple inheritance



# Fast Type Inclusion Tests – Binary Matrix

Class types

	C1	C2	C3	C4	C5
C1	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>
C2	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>
C3	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>
C4	<b>1</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>1</b>
C5	<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>

Class types

Tests are efficient, but Matrix will be large in practice. The matrix can be compacted, but this increases access time. This can handle multiple inheritance also.

**BM [C<sub>i</sub>, C<sub>j</sub>] = 1, iff C<sub>i</sub> is a subclass of C<sub>j</sub>**

# Relative (Schubert's) Numbering

$\{l_a, r_a\}$  for a node  $a$  :

$r_a$  is the ordinal number of the node  $a$  in a **postorder traversal** of the tree. Let  $\triangleleft$  denote “**subtype of**” relation. All descendants of a node are subtypes of that node.

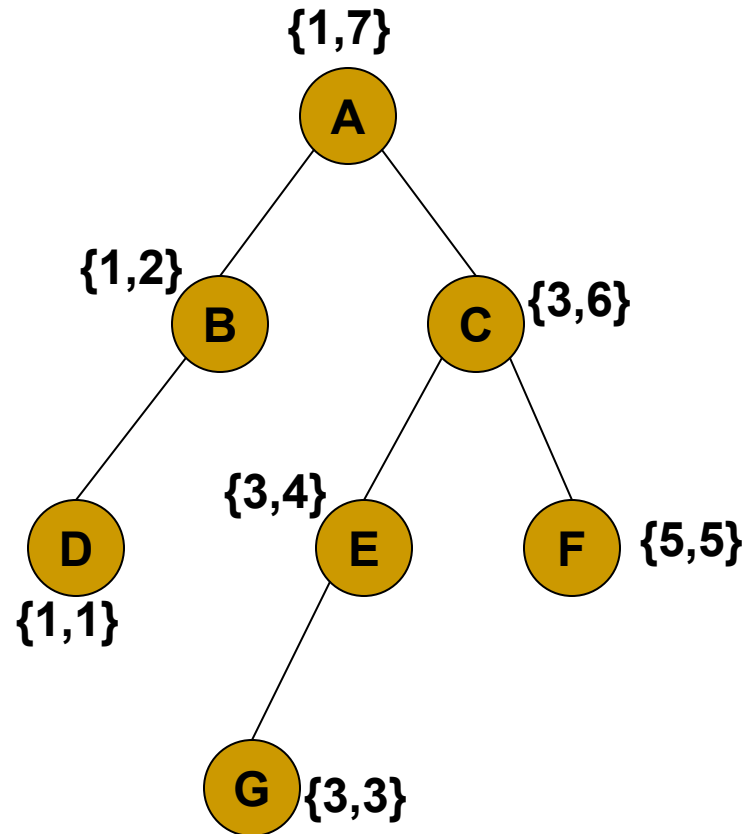
$\triangleleft$  is reflexive and transitive.

$l_a = \min \{ r_p \mid p \text{ is a descendant of } a \}$ .

Now,  $a \triangleleft b$ , iff  $l_b \leq r_a \leq r_b$ .

This test is very fast and is  $O(1)$ .  
Works only for single inheritance.

**Extensions to handle multiple inheritance are complex.**



# Devirtualization – Class Hierarchy Analysis

- Reduces the overhead of virtual method invocation
- Statically determines which virtual method calls resolve to **a single method**
- Such calls are either inlined or replaced by static calls
- Builds a class hierarchy and a call graph

# Class Hierarchy Analysis

```
class X extends object {  
    void f1() { . . . }  
    void f2() { . . . }  
}  
class Y extends X {  
    void f1() { . . . }  
}  
class Z extends X {  
    void f1() { . . . }  
    public static void main(...) {  
        X a = new X(); Y b = new Y();  
        Z c = new Z();  
        if (...) a = c;  
        // other code  
        a.f1(); b.f1(); b.f2();  
    }  
}
```

