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A graph G is called t-tough where t>0 is any real number, if for every separator S, the graph G-S has at most |S|/t components. Hamiltonian graphs are clearly 1-tough.

Toghness Conjecture (Chvátal 1973): There exists an integer t such that every t-tough graph has a Hamilton Cycle.

The 4-color conjecture can be reduced to simple 3-connected maximal planar graphs, i.e to 3-connected triangulations. Considering the dual, we get that 4-color conjecture is equivalent to the assertion that every 3-connected cubic plane graph is 4-face colorable.

A 3-connected cubic plane graph is 4-face colorable if and only if it is 3-edge colorable. Thus 4-color conjecture is equivalent to the assertion that every 3-connected cubic graph is 3-edge colorable.

If every 3-connected cubic graph is hamiltonian the above assertion is true. Tutte demonstrated that this is not true.

Every 4-connected planar graph has a hamiltonian cycle.