Numerical Optimization Unconstrained Optimization

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NPTEL Course on Numerical Optimization

Unconstrained Minimization Algorithm

\n- (1) Initialize
$$
x^0
$$
 and ϵ , set $k := 0$.
\n- (2) while $||g(x^k)|| > \epsilon$
\n- (a) Find a descent direction d^k for f at x^k
\n- (b) Find α^k (> 0) along d^k such that
\n- (i) $f(x^k + \alpha^k d^k) < f(x^k)$
\n- (ii) α^k satisfies Armijo-Wolfe conditions
\n- (c) $x^{k+1} = x^k + \alpha^k d^k$
\n- (d) $k := k + 1$
\n- **endwhile**
\n

Output : $x^* = x^k$, a stationary point of $f(x)$.

Does this algorithm converge?

Consider the problem,

$$
\min_{\mathbf{x}} f(\mathbf{x})
$$

- Let $f \in \mathcal{C}^1$ and *f* be bounded below.
- An optimization algorithm to minimize $f(x)$ generates a sequence, $\{x^k\}, k \geq 0$.
- Let the corresponding sequence of function values be ${f^k}, k \ge 0.$
- $f^{k+1} < f^k, \ k \geq 0$
- Stopping condition: $\|\boldsymbol{g}^k\| < \epsilon$

What can we say about $\|{\boldsymbol{g}}^{k}\|$ as $k\to\infty$?

Suppose, at every iteration *k* of the optimization algorithm,

- The direction \boldsymbol{d}^k is chosen such that $\boldsymbol{g}^{kT} \boldsymbol{d}^k < 0$
- Define $\phi(\alpha) = f(\pmb{x}^k + \alpha \pmb{d}^k)$. $\alpha^k (> 0)$ is chosen such that Armijo-Wolfe conditions are satisfied.

$$
f^{k+1} \leq f^k + c_1 \alpha \mathbf{g}^{k} \mathbf{g}^{k}, \ c_1 \in (0, 1) \n\phi'(\alpha^k) \geq c_2 \phi'(0), \ c_2 \in (c_1, 1)
$$

$$
\bullet \ f^{k+1} < f^k \ \forall \ k \ge 0
$$
\n
$$
\bullet \ x^{k+1} = x^k + \alpha^k d^k
$$

Given: $f^{k+1} < f^k \ \forall \ k \geq 0$.

 ${f^k}$: Monotonically decreasing sequence, which is also bounded below.

$$
\therefore \{f^k\} \to f^* \text{ where } f^* < \infty.
$$

$$
\therefore f^0 - f^k < \infty \ \forall \ k \ge 0
$$

$$
\therefore \lim_{k \to \infty} f^0 - f^k < \infty
$$

Using Armijo's condition, α^{j} 's are chosen such that

$$
f^{k+1} \leq f^k + c_1 \alpha^k \mathbf{g}^{k} \mathbf{d}^k
$$

$$
\leq f^0 + c_1 \sum_{j=0}^k \alpha^j \mathbf{g}^j \mathbf{d}^j
$$

Therefore,

$$
\infty > f^0 - f^{k+1} \ge -c_1 \sum_{j=0}^k \alpha^j \mathbf{g}^{jT} \mathbf{d}^j
$$

Therefore, sum of infinitely many positive terms is *finite*. This implies, beyond certain iteration *k*, $\alpha^k g^{kT} d^k = 0$. Using Wolfe condition, α^k is chosen such that

$$
\begin{array}{rcl}\phi'(\alpha^k) & \geq & c_2 \; \phi'(0), \; c_2 \in (c_1,1) \\[3mm] \therefore \, \, {\boldsymbol g}^{k+1^T} {\boldsymbol d}^k & \geq & c_2 \, \, {\boldsymbol g}^{k^T} {\boldsymbol d}^k \\[3mm] \therefore \, \, ({\boldsymbol g}^{k+1} - {\boldsymbol g}^k)^T {\boldsymbol d}^k & \geq & (c_2-1) {\boldsymbol g}^{k^T} {\boldsymbol d}^k\end{array}
$$

Let *g* be Lipschitz continuous. That is, $\exists L, 0 < L < \infty$ such that

$$
\|{\bm{g}}^{k+1}-{\bm{g}}^{k}\|\leq L\|{\bm{x}}^{k+1}-{\bm{x}}^{k}\|
$$

But, we have, $\mathbf{x}^{k+1} = \mathbf{x}^k + \alpha^k \mathbf{d}^k$.

$$
\therefore \|\mathbf{g}^{k+1} - \mathbf{g}^k\| \leq L\alpha^k \|\mathbf{d}^k\|
$$

$$
\therefore (\mathbf{g}^{k+1} - \mathbf{g}^k)^T \mathbf{d}^k \leq L\alpha^k \mathbf{d}^{k} \mathbf{d}^k
$$

But, using Wolfe conditions, $(g^{k+1} - g^k)^T d^k \ge (c_2 - 1)g^{kT} d^k$. Therefore,

$$
\alpha^k \geq \frac{c_2 - 1}{L} \frac{\mathbf{g}^{kT} \mathbf{d}^k}{\|\mathbf{d}^k\|^2}
$$

$$
\therefore \alpha^k \mathbf{g}^{kT} \mathbf{d}^k \leq \frac{c_2 - 1}{L} \frac{(\mathbf{g}^{kT} \mathbf{d}^k)^2}{\|\mathbf{d}^k\|^2}
$$

$$
\therefore -c_1 \alpha^k \mathbf{g}^{kT} \mathbf{d}^k \geq c_1 \frac{(1 - c_2)}{L} \frac{(\mathbf{g}^{kT} \mathbf{d}^k)^2}{\|\mathbf{d}^k\|^2}
$$

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$$
-c_1 \alpha^k {\bf g}^{kT} {\bf d}^k \ \ \ge \ \ c_1 \frac{(1-c_2)}{L} \frac{({\bf g}^{kT} {\bf d}^k)^2}{\|{\bf d}^k\|^2}
$$

Let θ_k be the angle between g^k and d^k . Therefore,

$$
-c_1 \alpha^k \mathbf{g}^{kT} \mathbf{d}^k \geq c_1 \frac{(1-c_2)}{L} \frac{\|\mathbf{g}^k\|^2 \|\mathbf{d}^k\|^2 \cos^2 \theta_k}{\|\mathbf{d}^k\|^2}
$$

$$
\therefore -c_1 \alpha^k \mathbf{g}^{kT} \mathbf{d}^k \geq c_1 \frac{(1-c_2)}{L} \|\mathbf{g}^k\|^2 \cos^2 \theta_k
$$

But, using Armijo's conditions, $-c_1 \sum_{k=0}^{\infty} \alpha^k \mathbf{g}^{k^T} \mathbf{d}^k < \infty$. Therefore,

$$
c_1\frac{(1-c_2)}{L}\sum_{k=0}^{\infty}||\mathbf{g}^k||^2\cos^2\theta_k < \infty
$$

$$
c_1 \frac{(1-c_2)}{L} \sum_{k=0}^{\infty} \|g^k\|^2 \cos^2 \theta_k < \infty
$$

This implies

$$
\|\mathbf{g}^k\|^2 \cos^2 \theta_k \to 0.
$$

If, at every iteration, d^k is chosen such that,

$$
\mathbf{g}^{k^T}\mathbf{d}^k < 0 \text{ and } \cos^2\theta_k \ge \delta > 0,
$$

then, we have,

$$
\lim_{k\to\infty}\|\boldsymbol{g}^k\|=0.
$$

Global Convergence Theorem

Global Convergence Theorem [*Zoutendijk*]

Consider the problem to minimize $f(x)$ over \mathbb{R}^n . Suppose f is bounded below in \mathbb{R}^n , $f \in \mathcal{C}^1$ and the gradient, $\nabla f(= g)$ is Lipschitz continuous. If at every iteration *k* of an optimization algorithm, a descent direction d^k is chosen such that $\cos^2 \theta_k > \delta(>0)$ (where θ_k is the angle between d^k and g^k) and α^k satisfies Armijo-Wolfe conditions, then the optimization algorithm either *terminates in a finite number of iterations* or

$$
\lim_{k\to\infty} \|{\boldsymbol{g}}^k\|=0.
$$

Sufficient Decrease and Backtracking

Armijo-Goldstein Conditions: Choose α^k such that

 $\phi_2(\alpha^k) \leq f(\pmb{x}^k + \alpha^k \pmb{d}^k) \leq \phi_1(\alpha^k)$

where
$$
\phi_1(\alpha) = f(x^k) + c_1 \alpha g^{kT} d^k
$$
, $c_1 \in (0, 1)$ and
\n $\phi_2(\alpha) = f(x^k) + c_2 \alpha g^{kT} d^k$, $c_2 \in (c_1, 1)$.

Use of *backtracking* line search with Armijo's condition

Backtracking Line Search

(1) Choose $\hat{\alpha}(> 0), \rho \in (0, 1), c_1 \in (0, 1)$. Set $\alpha = \hat{\alpha}$. (2) while $f(x^k + \alpha d^k) > f(x^k) + c_1 \alpha g^{kT} d^k$ $\alpha := \rho \alpha$ endwhile **Output :** $\alpha^k = \alpha$

Descent direction set: $\{ \boldsymbol{d} \in \mathbb{R}^n : \boldsymbol{g}^{\boldsymbol{k}^T} \boldsymbol{d} < 0 \}$ where $\boldsymbol{g}^k = \boldsymbol{g}(\boldsymbol{x}^k)$

Descent Directions

- Let $g^k \neq 0$ and $d^k = -A^k g^k$ where A^k is a symmetric matrix
- If A^k is positive definite,

$$
\mathbf{g}^{k^T} \mathbf{d}^k = -\mathbf{g}^{k^T} A^k \mathbf{g}^k < 0
$$

\n
$$
\Rightarrow \mathbf{d}^k \text{ is a descent direction}
$$

 $d^k = -A^k g^k$ is a *descent direction* if A^k is positive definite. Different optimization algorithms use different *A k*

How to find d^k ?

Consider the first order approximation to $f(x)$ about x^k :

$$
f(\mathbf{x}) \approx \hat{f}(\mathbf{x}) \stackrel{\Delta}{=} f(\mathbf{x}^k) + \mathbf{g}^{k^T}(\mathbf{x} - \mathbf{x}^k) = f(\mathbf{x}^k) + \mathbf{g}^{k^T} \mathbf{d}
$$

Maximum decrease in $\hat{f}(x)$ is possible by solving (P1):

$$
\begin{array}{ll}\n\min_{\boldsymbol{d}} & \boldsymbol{g}^{k^T} \boldsymbol{d} \\
\text{s.t.} & \boldsymbol{d}^T \boldsymbol{d} = 1\n\end{array}
$$

Let θ_k be the angle between g^k and d .

$$
\mathbf{g}^{k^T}\mathbf{d} = \|\mathbf{g}^k\| \|\mathbf{d}\| \cos \theta_k
$$

=
$$
\|\mathbf{g}^k\| \cos \theta_k \ (\because \mathbf{d}^T \mathbf{d} = 1)
$$

Therefore, the solution to the problem (P1) is $d^k = -g^k / \|g^k\|$

Steepest Descent Method

• Uses the steepest descent direction,
$$
d^k = -g^k
$$

Steepest Descent Algorithm

(1) Initialize x^0 and ϵ , set $k := 0$. (2) while $\|\mathbf{g}^k\| > \epsilon$ (a) $d^k = -g^k$ (b) Find α^k (> 0) along d^k such that (i) $f(x^k + \alpha^k d^k) < f(x^k)$ (ii) α^k satisfies Armijo-Wolfe conditions (c) $x^{k+1} = x^k + \alpha^k d^k$ (d) $k := k + 1$ endwhile

Output : $x^* = x^k$, a stationary point of $f(x)$.

• Exact or Backtracking line search can be used in step 2(b)

$$
\min \; f(x) \stackrel{\Delta}{=} (x_1 - 7)^2 + (x_2 - 2)^2
$$

Behaviour of the steepest descent algorithm (with exact line search) applied to $f(x)$ using $x^0 = (9, 4)^T$

Behaviour of the steepest descent algorithm (with exact line search) applied to $f(x)$ using $x^0 = (5.5, 3)^T$

$$
\min f(x) \stackrel{\Delta}{=} 4x_1^2 + x_2^2 - 2x_1x_2
$$

\n• $g(x) = \begin{pmatrix} 8x_1 - 2x_2 \\ 2x_2 - 2x_1 \end{pmatrix}, H(x) = \begin{pmatrix} 8 & -2 \\ -2 & 2 \end{pmatrix}.$
\n• $x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$

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$$
\min f(x) \stackrel{\Delta}{=} 4x_1^2 + x_2^2 - 2x_1x_2
$$

Behaviour of the steepest descent algorithm (with exact line search) applied to $f(x)$ using $x^0 = (-1, -2)^T$

$$
\min f(x) \stackrel{\Delta}{=} 4x_1^2 + x_2^2 - 2x_1x_2
$$

Behaviour of the steepest descent algorithm (with exact line search) applied to $f(\mathbf{x})$ using $\mathbf{x}^0 = (1,0)^T$

$$
\min f(\mathbf{x}) \stackrel{\Delta}{=} 100(x_2 - x_1^2)^2 + (1 - x_1)^2
$$
\n• $\mathbf{x}^* = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$

Behaviour of the steepest descent algorithm (with backtracking line search) applied to $f(x)$ using $x^0 = (0.6, 0.6)^T$

$$
\min f(\mathbf{x}) \stackrel{\Delta}{=} 100(x_2 - x_1^2)^2 + (1 - x_1)^2
$$

Table: Steepest descent method (with backtracking line search) applied to Rosenbrock function, using $x^0 = (0.6, 0.6)^T$.

Behaviour of the steepest descent algorithm (with backtracking line search) applied to $f(x)$ using $x^0 = (-1.2, 1)^T$

$$
\min f(\mathbf{x}) \stackrel{\Delta}{=} 100(x_2 - x_1^2)^2 + (1 - x_1)^2
$$

Table: Steepest descent method (with backtracking lines search) applied to Rosenbrock function, using $x^0 = (-1.2, 1)^T$.