

Numerical Differentiation

Note Title

3/23/2011

$$f'(a) = \frac{f(a+h) - f(a-h)}{2h} - \frac{h^2}{6} f'''(\xi)$$

$$f'_{\text{comp}} = \frac{f(a+h) + E_1 - \{f(a-h) + E_2\}}{2h} \quad \begin{array}{l} E_1, E_2 : \\ \text{round-off} \\ \text{errors} \end{array}$$

$$f'(a) = f'_{\text{comp}} + \frac{E_2 - E_1}{2h} - \frac{h^2}{6} f'''(\xi)$$

↓
does not decrease

Numerical Integration

$$\int_a^b f(x) dx \approx \frac{h}{2} (f(a) + f(b)) + h \sum_{i=1}^{n-1} f(t_i) = T_n .$$

$$T_n^{\text{comp.}} = T_n + \frac{h}{2} (E_0 + E_n) + h \sum_{i=1}^{n-1} E_i$$

Round-off error

System of linear equations

Note Title

2/7/2011

$$a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n = b_1$$

$$a_{21} x_1 + a_{22} x_2 + \dots + a_{2n} x_n = b_2$$

⋮

$$a_{n1} x_1 + a_{n2} x_2 + \dots + a_{nn} x_n = b_n$$

$Ax = b$, $A = [a_{ij}]$: coefficient matrix

$b = [b_1 \ b_2 \ \dots \ b_n]^T$: Right Hand Side

$x = [x_1 \ x_2 \ \dots \ x_n]^T$: unknown vector

$[A \ b]$: $n \times (n+1)$: augmented matrix.

$Ax = b$. If A is invertible, then the system has a unique solution: $x = A^{-1}b$.

Cramer's Rule

$$x_j = \frac{\begin{vmatrix} a_{11} & \dots & b_1 & a_{1n} \\ a_{21} & & b_2 & a_{2n} \\ \vdots & & \vdots & \vdots \\ a_{n1} & & b_n & a_{nn} \end{vmatrix}}{\det A}, \quad j = 1, \dots, n.$$

too expensive

Upper triangular system.

$$u_{11} x_1 + u_{12} x_2 + \dots + u_{1n} x_n = y_1$$

$$u_{22} x_2 + \dots + u_{2n} x_n = y_2$$

$$u_{n-1,n-1} x_{n-1} + u_{n-1,n} x_n = y_{n-1}$$

$$u_{nn} x_n = y_n$$

$$\det(U) = u_{11} u_{22} \dots u_{nn} \neq 0$$

$$\Rightarrow u_{ii} \neq 0$$

$$x_n = \frac{y_n}{u_{nn}}, \quad x_{n-1} = \frac{y_{n-1} - u_{n-1,n} x_n}{u_{n-1,n-1}}, \dots$$

Aim: To replace $Ax = b$ by an equivalent upper triangular system $Ux = y$ using elementary row transformations of the type

$$R_i \rightarrow R_i - \alpha R_j,$$

where $\alpha \in \mathbb{R}$, R_i : i th row of the augmented matrix.

Additional Assumption

$$\begin{array}{c} A_k \\ \left[\begin{array}{cccc|ccc} a_{11} & a_{12} & \cdots & a_{1k} & \cdots & a_{1n} & \\ a_{21} & a_{22} & \cdots & a_{2k} & \cdots & a_{2n} & \\ \vdots & & & & & & \\ \vdots & & & & & & \\ a_{k1} & a_{k2} & \cdots & a_{kk} & \cdots & a_{kn} & \\ \vdots & & & & & & \\ a_{n1} & a_{n2} & \cdots & a_{nk} & \cdots & a_{nn} & \end{array} \right] \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_k \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_k \\ \vdots \\ b_n \end{bmatrix} . \end{array}$$

A_k : principal leading submatrix.

Assumption : $\det(A_k) \neq 0, k=1, 2, \dots, n$

Gauss elimination method

1 flop (floating point operation)

= 1 addition / subtraction + 1 multiplication / division.

Step 1:

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ a_{21} & a_{22} & \dots & a_{2n} & b_2 \\ \vdots & \vdots & \dots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} & b_n \end{bmatrix} \rightarrow \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ 0 & a_{22}^{(1)} & \dots & a_{2n}^{(1)} & b_2^{(1)} \\ \vdots & \vdots & \dots & \vdots & \vdots \\ 0 & a_{n2}^{(1)} & \dots & a_{nn}^{(1)} & b_n^{(1)} \end{bmatrix}$$

$$\begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ a_{21} & a_{22} & \dots & a_{2n} & b_2 \\ \vdots & & & & \\ a_{n1} & a_{n2} & \dots & a_{nn} & b_n \end{bmatrix} \rightarrow \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ 0 & a_{22}^{(1)} & \dots & a_{2n}^{(1)} & b_2^{(1)} \\ \vdots & & & & \\ 0 & a_{n2}^{(1)} & \dots & a_{nn}^{(1)} & b_n^{(1)} \end{bmatrix}$$

$a_{11} \neq 0$ (why?), Define $m_{i1} = \frac{a_{i1}}{a_{11}}$, $i = 2, \dots, n$

$R_i \rightarrow R_i - m_{i1} R_1$, $i = 2, \dots, n$

$a_{ij}^{(1)} = a_{ij} - m_{i1} a_{1j}$, $b_i^{(1)} = b_i - m_{i1} b_1$, $i = 2, \dots, n$
 $j = 2, \dots, n$.

$(n-1)n$ mult. + $(n-1)n$ subtractions + $n-1$ divisions

$$A_2 = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \rightarrow \begin{bmatrix} a_{11} & a_{12} \\ 0 & a_{22}^{(1)} \end{bmatrix}$$

$$\det(A_2) = a_{11} a_{22}^{(1)} \Rightarrow a_{22}^{(1)} \neq 0$$

‡
0 by assumption

Perform similar operations on $(n-1) \times n$ matrix.

$$[A : b] \rightarrow [A^{(1)} : b^{(1)}] \rightarrow [A^{(2)} : b^{(2)}]$$

$$[A : b] \rightarrow [A^{(1)} : b^{(1)}]$$

$(n-1)n$ multi. + $(n-1)n$ subtractions + $n-1$ divisions

$$[A^{(1)} : b^{(1)}] \rightarrow [A^{(2)} : b^{(2)}]$$

$(n-2)(n-1)$ mult. + $(n-2)(n-1)$ sub. + $n-2$ divisions

Total number of operations : $[A : b] \rightarrow [U : y]$

$$[(n-1)^2 + (n-2)^2 + \dots + 1] + [(n-1) + (n-2) + \dots + 1] \text{ flops}$$

$$+ [(n-1) + (n-2) + \dots + 1] \text{ divisions}$$

$$= \frac{(n-1)n(2n-1)}{6} + \frac{(n-1)n}{2} \text{ flops} + \frac{(n-1)n}{2} \text{ divisions}$$

$O\left(\frac{n^3}{3}\right)$

Back Substitution

$$Ax = b \rightarrow Ux = y. \quad U = [u_{ij}]$$

$$u_{ij} \neq 0, \quad i=1, \dots, n, \quad u_{ij} = 0 \quad \text{if } i > j.$$

$$u_{ii} x_i + u_{i,i+1} x_{i+1} + \dots + u_{in} x_n = y_i : \text{ith eq}^n$$

$$x_i = \frac{y_i - \sum_{j=i+1}^n u_{ij} x_j}{u_{ii}}, \quad i = n, n-1, \dots, 1.$$

$$\text{Number of operations} = \sum_{i=1}^{n-1} (n-i) \text{ flops} + n \text{ divisions}$$

$$= \frac{(n-1)n}{2} \text{ flops} + n \text{ divisions} \quad O\left(\frac{n^2}{2}\right)$$

Definition: A is said to be positive-definite

if i) $A^T = A$ ii) $x^T A x > 0$ if $x \neq \bar{0}$

$$A u = \lambda u, u \neq \bar{0} \Rightarrow \lambda = \frac{u^T A u}{u^T u} > 0$$

eigenvalues of A are > 0 .

$\det(A) = \text{product of eigenvalues} > 0$

A_k : principal leading submatrix of A of order k .

Claim: A_k is positive definite

Proof: $A^T = A \Rightarrow a_{ij} = a_{ji}, i, j = 1, \dots, n$
 $\Rightarrow a_{ij} = a_{ji}, i, j = 1, \dots, k \Rightarrow A_k^T = A_k$

Claim: If $y \neq \bar{0}$, then $y^T A_k y > 0$

Proof:

Let $y = [y_1, y_2, \dots, y_k]^T$ be such that at least one $y_p \neq 0$. Define $x = [y, \underbrace{0, 0, \dots, 0}_{n-k}]^T$.

$$\begin{aligned} \text{Then } y^T A_k y &= \sum_{i=1}^k \sum_{j=1}^k a_{ij} y_i y_j \\ &= \sum_{i=1}^n \sum_{j=1}^n a_{ij} x_i x_j = x^T A x > 0 \end{aligned}$$

Thus A_k is positive definite and hence

$$\det(A_k) > 0.$$

$A = M^T M$, M : invertible.

$$A^T = A. \quad x^T A x = x^T M^T M x \\ = y^T y = \sum_{j=1}^n y_j^2$$

$$x \neq \bar{0} \Rightarrow Mx \neq \bar{0} \Rightarrow y^T y > 0.$$

A : positive definite.

Converse is true: Cholesky decomposition